## Aliens Arrive at Moon Base

Here is a video walkthrough of part a, and here is a video walkthrough of part b.
a) Suppose the aliens have received a landing code as a series of strings, but it's all scrambled up! Help them sort the strings so they can land on the moon base! More specifically, suppose we want to sort $N$ distinct Strings, each of length $W$, that are comprised of characters from an alien alphabet of size $N^{2}$. We can visualize the set of strings to sort as the image below.


What are the runtimes of the following sorting algorithms, in the best and worst case? If we use some sorting algorithm, say quicksort, as the subroutine in LSD/MSD sort, that means we use quicksort to sort each character. If two Strings are distinct, that means that at least one character differs. Finally, if the sorting algorithm provided would not work, select "Would Not Work" for the best and worst case. Your answers may involve both $W$ and $N$.

Hint: When we compare two strings, it takes constant time in the best case, and $W$ time in the worst case.
i) Merge sort


Not Work
Worst Case:

| $\bigcirc \Theta(1) \bigcirc \Theta(\log N)$ | $\bigcirc \Theta(N)$ | $\bigcirc \Theta(N \log N)$ | $\bigcirc \Theta\left(N^{2}\right)$ | $\left.\bigcirc \Theta\left(N^{2}\right) \log N\right)$ |
| :--- | :---: | :---: | :---: | :---: | :---: |
| $\bigcirc \Theta(\log W)$ | $\bigcirc \Theta(W)$ | $\bigcirc \Theta(W \log W)$ | $\bigcirc \Theta\left(W^{2}\right)$ | $\bigcirc \Theta\left(W^{2} \log W\right)$ |
| $\bigcirc \Theta(W \log N)$ | $\bigcirc \Theta(W N)$ | $\bigcirc \Theta(W N \log N)$ | $\bigcirc \Theta\left(W N^{2}\right)$ | $\bigcirc \Theta\left(W N^{2} \log N\right)$ |
| $\Theta(N \log W)$ | $\bigcirc \Theta(N W \log W)$ | $\bigcirc \Theta\left(N W^{2}\right)$ | $\bigcirc \Theta\left(N W^{2} \log W\right)$ |  |

$\bigcirc \Theta\left(N^{2} W^{2}\right) \quad \bigcirc$ Never terminates (infinite loop) $\bigcirc$ None of the above
Would Not Work

## Solution:

Best Case:


Would Not Work
Worst Case:


Would Not Work
ii) Insertion sort

## Best Case:

| $\bigcirc \Theta(1)$ | $\bigcirc$ | $\Theta(\log N)$ | $\bigcirc \Theta(N)$ | $\bigcirc \Theta(N \log N)$ | $\bigcirc \Theta\left(N^{2}\right)$ | $\left.\bigcirc \Theta\left(N^{2}\right) \log N\right)$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\bigcirc \Theta(\log W)$ | $\bigcirc \Theta(W)$ | $\bigcirc \Theta(W \log W)$ | $\bigcirc \Theta\left(W^{2}\right)$ | $\bigcirc \Theta\left(W^{2} \log W\right)$ |  |  |
| $\bigcirc \Theta(W \log N)$ | $\bigcirc \Theta(W N)$ | $\bigcirc \Theta(W N \log N)$ | $\bigcirc \Theta\left(W N^{2}\right)$ | $\bigcirc \Theta\left(W N^{2} \log N\right)$ |  |  |
| $\bigcirc \Theta(N \log W)$ | $\bigcirc \Theta(N W \log W)$ | $\bigcirc \Theta\left(N W^{2}\right)$ | $\bigcirc \Theta\left(N W^{2} \log W\right)$ |  |  |  |
| $\Theta \Theta\left(N^{2} W^{2}\right)$ | $\bigcirc$ Never terminates | (infinite loop) | $\bigcirc$ | None of the above |  |  |
| Would Not Work |  |  |  |  |  |  |

## Worst Case:

 Would Not Work

## Solution:

## Best Case:

 Would Not Work
Worst Case:


Would Not Work
iii) MSD sort with heap sort as the subroutine.

## Best Case:

| $\bigcirc \Theta(1) \bigcirc \Theta$ | $\left.\Theta(\log N) \bigcirc \Theta(N) \bigcirc \Theta(N \log N) \bigcirc \Theta\left(N^{2}\right) \bigcirc \Theta\left(N^{2}\right) \log N\right)$ |
| :---: | :---: |
| $\Theta(\log W)$ | $\bigcirc \Theta(W) \bigcirc \Theta(W l o g W) \bigcirc \Theta\left(W^{2}\right) \bigcirc \Theta\left(W^{2} \log W\right)$ |
| $\Theta(W \log N)$ | $\bigcirc \Theta(W N) \bigcirc \Theta(W N \log N) \bigcirc \Theta\left(W N^{2}\right) \bigcirc \Theta\left(W N^{2} \log N\right)$ |
| $\Theta(N \log W)$ | $\bigcirc \Theta(N W \log W) \bigcirc \Theta\left(N W^{2}\right) \bigcirc \Theta\left(N W^{2} \log W\right)$ |
| $\Theta\left(N^{2} W^{2}\right)$ | $\bigcirc$ Never terminates (infinite loop) $\bigcirc$ None of the above |
| Would Not Work |  |

## Worst Case:

| $\bigcirc \Theta(1)$ | $\bigcirc \Theta(\log N)$ | $\bigcirc \Theta(N)$ | $\bigcirc \Theta(N \log N)$ | $\bigcirc \Theta\left(N^{2}\right)$ | $\left.\bigcirc \Theta\left(N^{2}\right) \log N\right)$ |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: |
| $\bigcirc \Theta(\log W)$ | $\bigcirc \Theta(W)$ | $\bigcirc \Theta(W \log W)$ | $\bigcirc \Theta\left(W^{2}\right)$ | $\bigcirc \Theta\left(W^{2} \log W\right)$ |  |
| $\bigcirc \Theta(W \log N)$ | $\bigcirc \Theta(W N)$ | $\bigcirc \Theta(W N \log N)$ | $\bigcirc \Theta\left(W N^{2}\right)$ | $\bigcirc \Theta\left(W N^{2} \log N\right)$ |  |
| $\bigcirc \Theta(N \log W)$ | $\bigcirc \Theta(N W \log W)$ | $\bigcirc \Theta\left(N W^{2}\right)$ | $\bigcirc \Theta\left(N W^{2} \log W\right)$ |  |  |
| $\Theta\left(N^{2} W^{2}\right)$ | $\bigcirc$ Never terminates | (infinite loop) | $\bigcirc$ None of the above |  |  |

Would Not Work

## Solution:

## Best Case:



Would Not Work

## Worst Case:

$\left.\bigcirc \Theta(1) \bigcirc \Theta(\log N) \bigcirc \Theta(N) \bigcirc \Theta(N \log N) \bigcirc \Theta\left(N^{2}\right) \bigcirc \Theta\left(N^{2}\right) \log N\right)$
$\bigcirc \Theta(\log W) \bigcirc \Theta(W) \bigcirc \Theta(W \log W) \bigcirc \Theta\left(W^{2}\right) \bigcirc \Theta\left(W^{2} \log W\right)$
$\bigcirc(W \log N) \bigcirc \Theta(W N) \quad \sqrt{ } \Theta(W N \log N) \bigcirc \Theta\left(W N^{2}\right) \bigcirc \Theta\left(W N^{2} \log N\right)$
$\bigcirc \Theta(N \log W) \bigcirc \Theta(N W \log W) \bigcirc \Theta\left(N W^{2}\right) \quad \bigcirc \Theta\left(N W^{2} \log W\right)$
$\bigcirc\left(N^{2} W^{2}\right) \bigcirc$ Never terminates (infinite loop) $\bigcirc$ None of the above
Would Not Work
b) For the following two parts, we will be calling LSD sort with conventional LSD sort as the subroutine! Conventional LSD sort uses counting sort as the subroutine. Since conventional LSD sort requires that the objects it's comparing are comprised of digits of a fixed radix, or base, we will decompose each character
from the alien alphabet into a number of a certain radix. We can visualize the first pass of the sorting algorithm as shown in the image below.


For each part below, find only the worst case runtime of calling LSD sort with LSD sort as the subroutine, if we decompose each character into a number of the given radix.

Here are some helpful hints before you begin.

- If we decompose a character of size $M$ into a number of radix $B$, the number of digits in the resulting number is $\log _{B} M$.
- $\log _{B} B=1$.
- $\log _{B} M^{a}=a \times \log _{B} M$.
i) Radix 10 .



## Solution:



Would Not Work

## Explanation:

Before we get ahead of ourselves, let's remind ourselves what the runtime of LSD sort is: $W^{*}$ runtime of subroutine. Well, what is the runtime of our subroutine? Well, it's conventional LSD sort, so the runtime is $W^{\prime}\left(N^{\prime}+R^{\prime}\right)$ for some $W^{\prime}, N^{\prime}$, and $R^{\prime}$. $N^{\prime}$ is $N . R^{\prime}$ is the size of the radix, which is $10 . W^{\prime}$ is the length of the number, which is $\log _{10} N$, or $\log (N)$. As such, we calculate the runtime of the
subroutine as $\log (N) \times(N+10)=N \log (N)$, which gives us the final runtime of $W N \log (N)$.
ii) Radix N.

| $\bigcirc \Theta(1)$ | $\bigcirc(\log N)$ | $\bigcirc \Theta(N)$ | $\bigcirc \Theta(N \log N)$ | $\bigcirc \Theta\left(N^{2}\right)$ | $\left.\bigcirc \Theta\left(N^{2}\right) \log N\right)$ |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: |
| $\bigcirc \Theta(\log W)$ | $\bigcirc \Theta(W)$ | $\bigcirc \Theta(W \log W)$ | $\bigcirc \Theta\left(W^{2}\right)$ | $\bigcirc \Theta\left(W^{2} \log W\right)$ |  |
| $\Theta(W \log N)$ | $\bigcirc \Theta(W N)$ | $\bigcirc \Theta(W N \log N)$ | $\bigcirc \Theta\left(W N^{2}\right)$ | $\bigcirc \Theta\left(W N^{2} \log N\right)$ |  |
| $\bigcirc \Theta(N \log W)$ | $\bigcirc \Theta(N W \log W)$ | $\bigcirc \Theta\left(N W^{2}\right)$ | $\bigcirc$ | $\Theta\left(N W^{2} \log W\right)$ |  |
| $\bigcirc \Theta\left(N^{2} W^{2}\right)$ | $\bigcirc$ Never terminates | (infinite loop) | $\bigcirc$ None of the above |  |  | Would Not Work

## Solution:

| $\bigcirc \Theta(1) \bigcirc \Theta$ | $\left.\Theta(\log N) \bigcirc \Theta(N) \bigcirc \Theta(N \log N) \bigcirc \Theta\left(N^{2}\right) \bigcirc \Theta\left(N^{2}\right) \log N\right)$ |
| :---: | :---: |
| $\Theta(\log W)$ | $\bigcirc \Theta(W) \bigcirc \Theta(W \log W) \bigcirc \Theta\left(W^{2}\right) \bigcirc \Theta\left(W^{2} \log W\right)$ |
| $\bigcirc(W \log N)$ | $\sqrt{ } \Theta(W N) \bigcirc \Theta(W N \log N) \bigcirc \Theta\left(W N^{2}\right) \bigcirc \Theta\left(W N^{2} \log N\right)$ |
| $\Theta(N \log W)$ | $\bigcirc \Theta(N W \log W) \bigcirc \Theta\left(N W^{2}\right) \bigcirc \Theta\left(N W^{2} \log W\right)$ |
| $\Theta\left(N^{2} W^{2}\right)$ | $\bigcirc$ Never terminates (infinite loop) $\bigcirc$ None of the above |
| Would Not Wor |  |

Explanation: Mirroring the logic from the previous part, the only changes now are that $R^{\prime}$ is $N$, and $W^{\prime}$ is 2 since $\log _{N} N^{2}=2 \times \log _{N} N=2 \times 1=2$. We get the runtime of the subroutine as $2 \times(N+N)$ or $N$, and we get the final runtime of $W * N$ ! Best sorting algorithm yet :).

