

1 Conceptual Shortest Paths

[Here is a video walkthrough of the solutions.](#)

Answer the following questions regarding shortest path algorithms for a **weighted, undirected graph**. If the statement is true, provide an explanation. If the statement is false, provide a counterexample.

- (a) (T/F) If all edge weights are equal and positive, the breadth-first search starting from node A will return the shortest path from a node A to a target node B.

True. If all edges are equal in weight, then the shortest path from A to each node is proportional to the number of nodes on the path, so breadth first search will return the shortest path.

- (b) (T/F) If all edges have distinct weights, the shortest path between any two vertices is unique.

False. Consider a case of 3 nodes where AB is 3, AC is 5, and BC is 2. Here, the two possible paths from A to C both are of length 5. In general, paths with greater number of edges end up getting penalized more than paths with fewer edges.

- (c) (T/F) **Adding** a constant positive integer k to all edge weights will not affect any shortest path between two vertices.

False. Consider a case of 3 nodes A, B, and C where AB is 1, AC is 2.5 and BC is 1. Clearly, the best path from A to C is through B, with weight 2. However, if we add 1 to each edge weight, suddenly the path going through B will have weight 4, while the direct path is only 3.5.

- (d) (T/F) **Multiplying** a constant positive integer k to all edge weights will not affect any shortest path between two vertices.

True. Suppose we have arbitrary nodes u and v . Let's say the shortest path from u to v , before the multiplication by k , was of total weight w . This implies that every other path from u to v was of total weight greater than w . After multiplying each edge weight by k , the total weight of the shortest path becomes $w * k$ and the total weight of every other path becomes some number greater than $w * k$. Therefore, the original shortest path doesn't change.